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Learning to Play Mastermind Well Using the Anti-Mind with Feeback Algorithm

Jose Barahona da Fonseca

Department of Electrical Engineering and Computer Science Faculty of Sciences and Technology New University of Lisbon 2829-516 Caparica, Portugal, Email: jbfo@fct.unl.pt

Abstract—In a previous work we developed the Anti-Mind algorithm. The Anti-Mind program simulated a good player of the Mastermind game, discovering the secret code defined by the human operator (a sequence of four integers in the interval [0 5]) very quickly. Then we used the algorithm of Anti-Mind to help and correct a human operator trying to discover the secret code defined by the computer resulting in the Anti-Mind with Feedback algorithm. In this paper, we revisited this work and developed another faster implementation of the Anti-Mind with Feedback algorithm which has the drawback that it does not know the set of next good guesses, it just compares each guess with the previous moves and accepts it if it is coherent with all the previous moves. Nevertheless, we introduced an option to generate the set of good guesses, i.e., the guesses that are coherent with all the previous moves. This implementation allows generalizing the Mastermind game to more than four digits and more than six colours. We begin to define rigorously what we mean by a guess coherent with a previous move, next we define what is a good guess and, then, we enunciate five hypotheses about the Anti-Mind algorithm namely one that guarantees that if we always play a good guess we will find the code in a finite bounded number of guesses. We propose a strategy to play Mastermind with the maximization of repetions at the beginning of the game which reduces the cognitive overload to play well and validate it with the Anti-Mind with Feedback algorithm. Finally we compare the Anti-Mind algorithm with the Ant-Mind with maximization of repetitions of the guesses through intensive simulations and conclude that the original Anti-Mind algorithm has a better average performance in terms of the number of guesses to break the secret code.

Keywords-artificial intelligence; mastermind game; anti-mind algorithm; anti-mind with feedback algorithm

I. INTRODUCTION

It seems that the Mastermind game was invented by M. Meyerowitz in 1973 [1]. It is a two players game, the code maker and the code breaker, and the code breaker must find the secret code in at most 10 guesses. The secret code has four colours that may be of one of six colours. In each move the code breaker makes a guess of four colours and the code maker answers with the number of white pegs, *cpe*, the number of correct digits in wrong position, and the number of black pegs, the number of correct digits in right position, *cpc*. In this work the six colours will be represented by the integers in the interval [0 5]. Our algorithm was named Anti-Mind because it was originally created to break the code created by the human player. Then we created a variation of this algorithm where the human plays as the code breaker and the Anti-Mind algorithm is used to generate in each guess the set of good guesses, i.e.,

the guesses coherent with the previous moves. If the guess does not belong to the set of good moves, then, it is considered a bad move and the user is asked if he wants to see the set of good moves and, then, he must make another guess [2]. In this paper, we show how the human player can improve his skills using the Anti-Mind with Feedback algorithm. When it remains only one hypothesis that must be the secret code and the user is asked to guess the code without any more information. From the use of this algorithm resulted a simplified faster algorithm and a strategy of playing Mastermind.

Our simulations point to a worst case performance of the Anti-Mind algorithm of 9 guesses and Donald E. Knuth [3] showed through exhaustive simulations that his strategy guarantees a maximum of 5 guesses. He showed that the best first guess to guarantee this worst case performance is 1122. After the work of Knuth many proposals were published, e.g. [4]-[7], but no one beat the worst case performance of Knuths strategy of 5 guesses, nevertheless some works improved the average performance over all possible secret codes. Nevertheless, our algorithm is much faster than previous algorithms, since at each stage it selects randomly the next guess from the set of good guesses that are coherent with previous moves. In this sense, the Anti-Mind algorithm can be considered a stochastic algorithm. On the contrary Knuth's algorithm for each good guess generates the next set of good moves for each of the 15 possible combinations of cpc and cpe and chooses the good guess that minimizes the maximum number of next good guesses [3].

This paper is organized as follows. In Section 2, we describe how a good guess is defined based on the concept of the *coherence* of a guess with a previous *move*. In Section 3, we describe the Anti-Mind algorithm. In Section 4, we describe the Anti-Mind with Feedback algorithm. In Section 5, we show how the Anti-Mind with Feedback algorithm can be used to learn to play mastermind well, in Section 6, we present a faster version of the Anti-Mind with feedback algorithm, in Section 7, we present exhaustive simulation results of the Anti-Mind algorithm and the Anti-Mind algorithm with maximization of the number of repetitions of the guesses and in Section 8, we present the conclusions and possible vectors of evolution of our work.

II. DEFINITION OF A MASTERMIND GOOD GUESS

The main idea behind the Anti-Mind algorithm is the *coherence* between a guess and a previous *move*. In Definition 1, we define what is a *move*.

Definition 1: A *move* is the triplet (*guess*, *cpc*, *cpe*) where *cpc* and *cpe* result from the comparison between the *guess* and the secret code.

In Definition 2, we define what is the *coherence* between a *guess* and a previous *move*.

Definition 2: Consider a previous move, move, with cpc correct digits in right position and cpe correct digits in wrong position and a guess, guess. Considering that the comparison between the guess and the move resulted in cpc_i and cpe_i , then, the guess is coherent with move if (1) is true.

$$cpc = cpc_i \text{ AND } cpe = cpe_i$$
 (1)

When a guess is coherent with all previous moves, then, we say that it is a *good guess*. On the contrary if the guess is not coherent with at least one previous move, then, we say that it is a *bad guess*. In this sense, we can say that *playing Mastermind well* happens when all guesses are *good guesses*.

So, we can define rigorously the set of good guesses as all combinations that are coherent with all previous moves. This is expressed by (2).

$$set_good_guesses = \{\forall combination : (2) \\ \forall move. cpc = cpc \ i \ AND \ cpe = cpe \ i\}$$

Once obtained the set of good guesses, we can verify if a given guess is a good guess, testing if it belongs to the set of good guesses. Alternatively, we can compare the *guess* with all previous moves, and if it is coherent with all of them, then, it is a *good guess*.

III. THE ANTI-MIND ALGORITHM

After each guess, the Anti-Mind algorithm obtains the new set of good guesses and selects randomly one of them as the next guess. This is much less computationally expensive than Knuth's algorithm but the Anti-Mind algorithm has a greater worst performance of 8 guesses. In this sense, we can say that the Anti-Mind algorithm is a stochastic algorithm. So we can enunciate five hypotheses that we will show in a near future work that characterize the behaviour of the Anti-Mind algorithm based on emperical data like intensive simulations.

Hypothesis 1: The Anti-Mind algorithm always finds the secret code in less than 9 guesses if the codemaker did not make any error in *cpc* and *cpe* for all previous moves.

Hypothesis 2: When the Anti-Mind algorithm finds the secret code in 8 guesses, if the codemaker did not make any error in *cpc* and *cpe* for all previous moves, the eigth *good* guess is always the secret code. In appendix we show some *enough information* games with 8 guesses.

Hypothesis 3: If the codemaker makes errors in *cpc* and/or *cpe* in at least one previous move, then, the Anti-Mind algorithm always reaches a situation of an empty set of good guesses in less than 9 guesses.

Hypothesis 4: If the game reaches a situation where there is only one good guess, then, this good guess must be the secret

code, if the codemaker did not make errors in the previous moves.

Hypothesis 5: Obtaining the new set of good guesses is equivalent to obtain all combinations/guesses that are coherent with all previous moves

The result of hypothesis 2 is the main idea behind our previous work of the Anti-Mind with an unlimited number of lies [8]. When it reaches the conclusion that there is at least one lie in previous moves, reaching the situation of an empty set of good guesses, then, it begins to test the hypothesis of one lie, removing one previous move from the set of previous moves and generating each time the set of good guesses. If it reaches a point where all manners of removing a previous move resulted in an empty set of good guesses, then, it begins to remove two previous moves and so on until it reaches a point where the set of good guesses has a cardinal 1. Then the removed previous moves are the moves with lies and the good guess must be the secret code [8].

Hypothesis 1 can be proved by exhaustive computer search, where for each possible secret code are generated all possible Mastermind games with good guesses and saved the number of guesses when the secret code is found. Since there are only 1296 possible Mastermind secret codes, this computer search is not prohibitive in terms of runtime.

The result of hypothesis 4 is the main idea behind the second version of the Anti-Mind with feedback algorithm where we do not have all good guesses in memory, and generate them whenever it is asked, comparing all combinations with the previous moves and printing the guesses that are coherent with all previous moves- see section 6.

IV. FIRST VERSION OF THE ANTI-MIND WITH FEEDBACK ALGORITHM

In the Anti-Mind with feedback algorithm the human is the codebreaker and the computer the codemaker. At each move, the new set of good guesses is generated, and when the guess does not belong to this set, the user is asked to enter another guess. There is an option that allows the user to see all the good guesses. If it is reached a situation where only one good guess remains, the user is asked to enter the secret code without any more information, since hypothesis 3 guarantees that the remaining good guess must be the secret code. In Table I we present a game of this type.

In Algorithm 1, we describe the Anti-Mind algorithm in detail.

V. USING THE ANTI-MIND WITH FEEDACK ALGORITHM TO FIND A GOOD PLAYING STRATEGY

Since we think in terms of symbolic expressions, and taking into account our cognitive limitations, it is easier to begin with a guess with repetitions like 0111. Now if the answer is cpc=3, cpe=0, we can conclude that there exist three 1s in the last three positions, OR one 0 in the first position AND two 1s in two of the last three positions. This symbolical logical expression corresponds to the set of 20 good guesses presented in Table II.

But if for the same secret code, the first guess was 0123, the answer would be cpc=2, cpe=0, which means a much more

Algorithm 1 Anti-Mind with feedback algorithm

```
n\_digits \leftarrowinput('Number of Digits=')
dig max \leftarrow input('Maximum Digit=')
n\_good\_guesses \Leftarrow (dig\_max+1)^n\_digits
secret\_code \leftarrow generate\_s\_c(n\_digits, dig\_max)
set\_good\_guesses \Leftarrow gen\_all\_gs(n\_digits, dig\_max)
counter \Leftarrow 0
while n_{good}_{guesses} > 1 do
  counter \Leftarrow counter + 1
  guess \leftarrow input('Guess =')
  cpc \leftarrow calc\_cpc(secret\_code, guess)
  cpe \leftarrow calc\_cpe(secret\_code, guess)
  if cpc == n_digits then
     display('You Found It')
     break
  end if
  flag\_bel \Leftarrow see\_if\_bel(guess, set\_good\_guesses)
  if flag_bel then
     display_cpc_cpe(cpc,cpe)
     set_good_guesses
     calc_new_set(guess, cpc, cpe, set_good_guesses)
  else
     display('Bad Move')
     in=input('Want to See Good Guesses?')
     if in == 1 then
       display_g_g(set_good_guesses)
     end if
  end if
end while
if n_{qood_{quesses}} = 1 then
  flaq \Leftarrow 0
  while 1 - flag do
     guess=input('Secret Code=')
     flaq \leftarrow (quess == secret \ code)
  end while
end if
```

complex symbolic expression: 0 and 1 are in right position OR 0 and 2 are in right position OR 0 and 3 are in right position OR 1 and 2 are in right position OR 1 and 3 are in right position OR 2 and 3 are in right position. This complex symbolical logical expression corresponds to a much greater set of 96 good guesses. In Table III we show these good guesses.

So in terms of *cognitive overload* it seams better to play in the beginning with repetitions.

VI. SECOND VERSION OF THE ANTI-MIND WITH FEEDBACK ALGORITHM

In the first version of the algorithm we first generate all combinations, and for each guess and cpc and cpe we generate a new set of good guesses, and we decide if the new guess is good, testing if the guess belongs to the set of good guesses. For a genarilized version of Mastermind with more digits and a greater maximum digit, this can take a lot of time in the first moves. So, inspired in this strategy, we created a new version of the algorithm where we only compare the guess with previous moves, and accept it if it is coherent with all previous moves. This algorithm does not have the information of the

TABLE I. EXAMPLE OF A GAME WITH 8 GUESSES

anti_mind_real(4,5,1, 1) Number of Possible Good Guesses=1296 move 1=0111 cpc=0 cpe=1 Number of Possible Good Guesses=308 move 2=1222 cpc=1 cpe=0 Number of Possible Good Guesses=90 move 3=1333 cpc=1 cpe=0 Number of Possible Good Guesses=20 move 4=1444 cpc=0 cpe=1 Number of Possible Good Guesses=6 move 5=4320 cpc=2 cpe=2 Number of Possible Good Guesses=3 move 6=4302 cpc=1 cpe=3 Number of Possible Good Guesses=2 move 7=4023 cpc=1 cpe=3 *ENOUGH INFORMATION** Secret Code=4230 *You Found It! in 8 Guesses, with 0 bad Guesses and 0 hints **

TABLE II. SET OF GOOD GUESSES AFTER GUESS 0111

0011	0101	0110	0112	0113	0114
0115	0121	0131	0141	0151	0211
0311	0411	0511	$1 \ 1 \ 1 \ 1 \ 1$	2111	3111
4111	5111				

TABLE III. SET OF GOOD GUESSES AFTER GUESS 0123

0003	0020	0022	0024	0025	0033
0043	0053	0100	0101	0104	0105
0110	0111	0114	0115	0140	0141
0144	0145	0150	0151	0154	0155
0220	0222	0224	0225	0303	0333
0343	0353	0403	0420	0422	0424
0425	0433	0443	0453	0503	0520
0522	0524	0525	0533	0543	0553
1113	1121	1 1 2 2	1124	1125	1133
1143	1153	2121	$2\ 1\ 2\ 2$	2124	2125
2223	2323	2423	2523	3113	3133
3143	3153	3223	3323	3423	3523
4113	4121	4122	4124	4125	4133
4143	4153	4223	4323	4423	4523
5113	5121	5122	5124	5125	5133
5143	5153	5223	5323	5423	5523

number of good guesses and it is impossibile the detection of an *enough information* situation, but it is much faster and we can play more difficult generalized Mastermind games inside Matlab environment. In Algorithm 2, we describe in detail this second version of the Anti-Mind with feedback algorithm.

VII. SIMULATION RESULTS

To evaluate the performance of the Anti-Mind algorithm, we made an exhaustive simulation over all possible secret codes, with 1000 runs for each secret code. In this simulation, the computer is the codemaker and the codebreaker. In Figure 1 we show the distribution of runs by the number of guesses to find the secret code. Then we repeat the simulation maximizing the number of repetitions of each guess. In Figure 2 we show

Algorithm 2 Second version of Anti-Mind with feedback

 $flag_not_found \leftarrow 1$ $n \ digits \leftarrow input('Number of Digits=')$ $dig_max \leftarrow input('Maximum Digit=')$ $secret_code \leftarrow generate_s_c(n_digits, dig_max)$ counter $\Leftarrow 0$ while *flag_not_found* do $counter \Leftarrow counter + 1$ $guess \leftarrow input('Guess =')$ $cpc \Leftarrow calc \ cpc(secret \ code, quess)$ $cpe \leftarrow calc_cpe(secret_code, guess)$ $cpc_o(counter) \Leftarrow cpc$ $cpe_o(counter) \Leftarrow cpe$ $guess_o(counter) \Leftarrow guess$ if $cpc = n_digits$ then display('You Found It') break end if for i=1:counter-1 do if $1 - flag_bad_guess_o(i)$ then $cpc_i \leftarrow calc_cpc(guess, guess_o(i))$ $cpe_i \leftarrow calc_cpe(guess, guess_o(i))$ $flag_bad_guess \iff 1 - (cpc_i = cpc) + 1 -$ $(cpe_i = cpe)$ if flag_bad_guess then break end if end if end for $flag_bad_guess_o(conter) \Leftarrow flag_bad_guess$ if flag bad guess then display('Bad Guess!') *in*=input('Do you want to see the good guesses?') if in == 1 then combination=zeros(1,n_digits) flag end $\Leftarrow 0$ while 1-flag_end do for i=1:counter-1 do $cpc_i \leftarrow calc_cpc(combination, guess_o(i))$ $cpe_i \leftarrow calc_cpe(combination, guess_o(i))$ $flag_coher \leftarrow 1 - (cpc_i = cpc_o(i)) + 1 (cpe_i = cpe_o(i))$ if flag coher==0 then break end if end for if flag_coher then display(combination) end if combination $\Leftarrow gen_n ext \ comb(combination)$ $flag_end \Leftarrow all_combs(combination)$ end while end if end if $flag_not_found \leftarrow 1 - (cpc = n_digits)$ end while

the results of this simulation. Comparing the two figures, we can say that the results of the second simulation are worse than the results of the first simulation, since we have a greater

percentage of runs with 5 guesses. This way we can say that playing mastermind *well*, with repetitions, has a lower performance than the Anti-Mind algorithm performance. So we confirm that the computer *thinks* better than the human [2].



FIGURE 1. NUMBER OF GUESSES IN PERCENTAGES DISTRIBUTION OF THE ANTI-MIND ALGORITHM.



FIGURE 2. NUMBER OF GUESSES IN PERCENTAGES DISTRIBUTION OF THE ANTI-MIND ALGORITHM WITH MAXIMIZATION OF REPETITIONS OF GUESSES.

VIII. CONCLUSIONS AND FUTURE WORK

We showed how to define rigorously what it is meant by playing Mastermind well and how the Anti-Mind algorithm can be used to learn playing well, resulting in the Anti-Mind with feedback algorithm. In the near future, we plan to demonstrate the five hypotheses enunciated in this work as theorems.

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APPENDIX

Enough Information Games with 8 Guesses

>> anti_mind_real(4,5,1, 1) Number of Possible Good Guesses=1296 move 1=0111 cpc=0 cpe=1 Number of Possible Good Guesses=308 move 2=1222 cpc=1 cpe=0 Number of Possible Good Guesses=90 move 3=1333 cpc=1 cpe=0 Number of Possible Good Guesses=20 move 4=1444 cpc=0 cpe=1 Number of Possible Good Guesses=6 move 5=4320 cpc=2 cpe=2 Number of Possible Good Guesses=3 move 6=4302 cpc=1 cpe=3 Number of Possible Good Guesses=2 move 7=4023 cpc=1 cpe=3 **ENOUGH INFORMATION** move 8=4230 **You Found It! in 8 Guesses, with 0 bad Guesses and 0 hints ** Number of Possible Good Guesses=1296 move 1=2201 cpc=0 cpe=2 Number of Possible Good Guesses=222 move 2=5320 cpc=3 cpe=0 Number of Possible Good Guesses=8 move 3=5322 cpc=2 cpe=0 Number of Possible Good Guesses=7 move 4=5520 cpc=2 cpe=0 Number of Possible Good Guesses=4 move 5=5310 cpc=2 cpe=0 Number of Possible Good Guesses=3 move 6=4320 cpc=3 cpe=0 Number of Possible Good Guesses=2 move 7=0320 cpc=3 cpe=0 **ENOUGH INFORMATION** move 8=3320 **You Found It! in 8 Guesses, with 0 bad Guesses and 0 hints ** Number of Possible Good Guesses=1296 move 1=2241 cpc=1 cpe=1 Number of Possible Good Guesses=230 move 2=2532 cpc=0 cpe=2 Number of Possible Good Guesses=34 move 3=5121 cpc=0 cpe=2 Number of Possible Good Guesses=10 move 4=1345 cpc=0 cpe=2 Number of Possible Good Guesses=6 move 5=4254 cpc=1 cpe=0 Number of Possible Good Guesses=3 move 6=3213 cpc=3

Number of Possible Good Guesses=2 move 7=3210 cpc=2 cpe=2 **ENOUGH INFORMATION** move 8=0213 **You Found It! in 8 Guesses, with 0 bad Guesses and 0 hints ** Number of Possible Good Guesses=1296 move 1=0214 cpc=1 cpe=0 Number of Possible Good Guesses=108 move 2=1111 cpc=0 cpe=0 Number of Possible Good Guesses=81 move 3=0300 cpc=0 cpe=1 Number of Possible Good Guesses=29 move 4=3444 cpc=1 cpe=0 Number of Possible Good Guesses=9 move 5=3232 cpc=3 cpe=0 Number of Possible Good Guesses=4 move 6=3233 cpc=2 cpe=0 Number of Possible Good Guesses=2 move 7=3252 cpc=3 cpe=0 **ENOUGH INFORMATION** move 8=3222 **You Found It! in 8 Guesses, with 0 bad Guesses and 0 hints ** Number of Possible Good Guesses=1296 move 1=1243 cpc=0 cpe=2 Number of Possible Good Guesses=312 move 2=2551 cpc=1 cpe=0 Number of Possible Good Guesses=50 move 3=3450 0=202 cpc=0 cpc=2 Number of Possible Good Guesses=12 move 4=4001 cpc=1 Number of Possible Good Guesses=4 move 5=2024 cpc=0 cpe=1 Number of Possible Good Guesses=3 move 6=0331 cpc=3 cpe=0 Number of Possible Good Guesses=2 move 7=0131 cpc=2 cpe=2 **ENOUGH INFORMATION** move 8=0311 **You Found It! in 8 Guesses, with 0 bad Guesses and 0 hints **